

# Bilevel Programming and MPCC : Connections and Counterexamples

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# Bilevel Programming Problem

$$\min_x F(x, y), \quad \text{subject to } x \in X, y \in S(x),$$

where for each  $x \in X$  the set  $S(x)$  is given as

$$S(x) = \operatorname{argmin}_y \{f(x, y) : y \in K(x)\},$$

where  $F : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$ ,  $f : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$  and  $K(x)$  is a closed convex set in  $\mathbb{R}^m$  depending on  $x \in X$ .

In our presentation we shall restrict ourselves to the case where  $X = \mathbb{R}^n$  in most situations. The set  $K(x)$  will often be given as

$$K(x) = \{y \in \mathbb{R}^m : g_i(x, y) \leq 0, i = 1, \dots, k\},$$

here  $y \mapsto g_i(x, y)$  is convex for each  $i$ .

# Optimistic and Pessimistic Formulation

The optimistic formulation is given as follows : Consider that  $S(x) \neq \emptyset$  for each  $x$  and define the function

$$\varphi_0(x) = \min_{y \in S(x)} F(x, y).$$

Then the *optimistic problem* is to minimize  $\varphi_0$  over  $x$ . We shall refer to the optimistic problems as  $(BP_o)$ .

The pessimistic formulation is given as follows : Let us define the function

$$\varphi_p(x) = \max_{y \in S(x)} F(x, y).$$

Thus the pessimistic bilevel problem consist of minimizing  $\varphi_p$  over  $\mathbb{R}^n$ . Note that the pessimistic formulation of a bilevel problem is viewed as one where the follower does not cooperate with the leader.

# Optimistic Bilevel Programming

This optimistic bilevel programming problem which is denoted as (OBP) is given as

$$\min_{x,y} F(x,y), \quad \text{subject to } y \in S(x).$$

Most researchers speak of this formulation as the bilevel programming. How is this problem related to the original optimistic formulation. How is  $(BP_o)$  is related to (OBP).

# Relation between optimistic formulation and OBP

Result 1 :

Let  $\bar{x}$  be the local solution of the optimistic formulation ( $BP_o$ ). Then for any  $\bar{y} \in S(\bar{x})$ , the vector  $(\bar{x}, \bar{y})$  is a local minimum of (OBP) if  $\bar{y}$  be such that  $\varphi_o(\bar{x}) = F(\bar{x}, \bar{y})$ .

Result 2 :

Let  $(\bar{x}, \bar{y})$  be the global minimizer of (OBP). Then  $\bar{x}$  is a global minimizer of the problem ( $BP_o$ ).

Result 3 : Let  $(\bar{x}, \bar{y})$  be the global minimizer of (OBP). Then  $\bar{x}$  is a global minimizer of the problem ( $BP_o$ ).

# KKT reformulation of OBP

Let the set  $K(x)$  be defined by convex inequalities. The KKT reformulation of (OBP) is given below and is called (OBP-KKT)

$$\min_{x,y} F(x,y), \quad \text{subject to} \quad x \in X, \nabla L(x,y,u) = 0, \quad u \geq 0, u^T g(x,y) = 0,$$

where  $L(x,y,u)$  is the Lagrangian function associated with the lower-level problem. The set of Lagrangian multipliers of the lower-level problem is given as

$$\Lambda(x,y) = \{u : \nabla L(x,y,u) = 0, \quad u \geq 0, u^T g(x,y) = 0\}$$

The set set  $X$  is often  $\mathbb{R}^n$ .

# The Global Case

Result 4 : Let  $(\bar{x}, \bar{y})$  be a global minimizer of (OBP) and the Slater CQ holds for the lower-level problem at  $x = \bar{x}$ . Then for any  $\bar{u} \in \Lambda(\bar{x}, \bar{y})$ , we have that  $(\bar{x}, \bar{y}, \bar{u})$  is a solution of (OBP-KKT).

Result 5 : Let  $(\bar{x}, \bar{y}, \bar{u})$  be the global minimizer of (OBP-KKT) . Let us assume that the Slater constraint qualification holds true for the lower-level problem for each  $x \in X$  . Then  $(\bar{x}, \bar{y})$  solves (OBP).

## Example : Global Case

Consider the following (OBP) in  $\mathbb{R}^2$ .

$$\min_{x,y} (x-1)^2 + y^2, \quad x \in \mathbb{R}, y \in S(x),$$

$$S(x) = \operatorname{argmin}_y \{x^2 y : y^2 \leq 0\}$$

Solution of (SBP) :  $(\bar{x}, \bar{y}) = (0, 0)$ .

Associated MPCC problem

$$\min_{x,y,\lambda} (x-1)^2 + y^2; \quad \text{subject to } x^2 + 2\lambda y = 0, \quad \lambda \geq 0, \quad y^2 \leq 0, \lambda y^2 = 0.$$

All feasible points of the MPCC is of the form  $(0, 0, \lambda)$  thus by solving the MPCC we cannot solve the bilevel problem in the context of global minimizers.



## Example : Local case

Consider the following (OBP) in the  $\mathbb{R}^2$ .

$$\min_{x,y} (x-1)^2 + (y-1)^2, \quad \text{subject to } x \in \mathbb{R}, y \in S(x),$$

where

$$S(x) = \operatorname{argmin}_y \{-y : x + y \leq 1, -x + y \leq 1\}.$$

The problem (OBP) has a unique global minimizer  $(\bar{x}, \bar{y}) = (0.5, 0.5)$  and there are no local minimizers.

## Example : Contd

The corresponding MPCC is given as

$$\begin{aligned} & \min_{x,y,\lambda} (x-1)^2 + (y-1)^2, \\ & \text{subject to} \\ & -1 + u_1 + u_2 = 0, \quad u_1 \geq 0, \quad u_2 \geq 0 \\ & u_1(x+y-1) = 0 \\ & u_2(x+y-1) = 0 \\ & x+y-1 \leq 0 \\ & -x+y-1 \leq 0. \end{aligned}$$

For example  $(x^*, y^*, u_1^*, u_2^*) = (0, 1, 0, 1)$  is a local solution of MPCC but  $(0, 1)$  is not a local solution of (OBP).

# Main Result : Local Case

Let  $\bar{x}$  be a point where Slater condition holds for the lower-level problem.  
Let  $\bar{y}$  be a solution of the lower-level problem corresponding to  $\bar{x}$ .  
For each  $\bar{u} \in \Lambda(\bar{x}, \bar{y})$  the point  $(\bar{x}, \bar{y}, \bar{u})$   
is a local minimizer of (OBP).  
Then  $(\bar{x}, \bar{y})$  is a local minimizer of (OBP).

# Simple Bilevel Programming Problem

Let us consider the following Simple Bilevel Programming (SBP) problem

$$\begin{aligned} & \text{minimize } f(x) \\ & \text{subject to} \\ & x \in \operatorname{argmin}\{h(x) : g(x) \leq 0\}. \end{aligned} \tag{1}$$

# KKT conditions

Now the question is if the lower level problem i.e.

$$\begin{aligned} & \text{minimize } h(x) \\ & \text{subject to} \\ & \quad g(x) \leq 0. \end{aligned} \tag{2}$$

can be replaced by its Karush Kuhn Tucker conditions?

The answer is yes if the Slater's CQ condition holds for the lower level problem.

# Simple MPCC problem

If the lower level problem of the SBP (1) is replaced by the KKT conditions then we get the following simple MPCC problem

$$\begin{aligned} & \text{minimize } f(x) \\ & \text{subject to} \\ & \quad \nabla h(x) + \lambda^t \nabla g(x) = 0 \\ & \quad g(x) \leq 0 \\ & \quad \lambda \geq 0 \\ & \quad \lambda^t g(x) = 0. \end{aligned} \tag{3}$$

For any  $x \in \mathbb{R}^n$  such that  $g(x) \leq 0$ , Let us define

$$\Lambda(x) := \{\lambda \geq 0 : \nabla h(x) + \lambda^t \nabla g(x) = 0, \lambda^t g(x) = 0\}$$

Then  $(x, \lambda)$  is a feasible point of the problem (2).

### Theorem

*Let  $\bar{x}$  is a global optimal solution of the simple bilevel programming problem and assume that the lower level problem satisfies the Slater's CQ condition i.e.  $\exists x \in \mathbb{R}^n$  such that  $g(x) < 0$ . Then for any  $\lambda \in \Lambda(\bar{x})$ , the point  $(\bar{x}, \lambda)$  is a global optimal solution of the corresponding MPCC Problem.*



## Theorem

*Let the Slater's condition holds for the lower level problem (2) of the SBP . Then  $(\bar{x}, \bar{\lambda})$  is a local solution of the MPCC problem implies that  $\bar{x}$  is a global solution of the SBP problem.*

## Corollary

*Let the Slater's condition holds for the lower level problem (2) of the SBP . Then  $(\bar{x}, \bar{\lambda})$  is a local solution of the corresponding MPCC problem implies that  $(\bar{x}, \bar{\lambda})$  is a global solution of the same.*

### Example

Slater's condition holds and the solution of SBP and MPCC are same.

Let

$$f(x) = \left(x - \frac{1}{2}\right)^2$$

$$h(x) = \begin{cases} 0 & \text{if } 0 \leq x \leq 1 \\ 1 & \text{if } x > 1 \end{cases}$$

$$g_1(x) = -x$$

$$g_2(x) = x - 3$$

Then Slater's condition holds as  $g_1(2) < 0$  and  $g_2(2) < 0$ .

Here the feasible set for the MPCC problem is

$$\{(x, \lambda_1, \lambda_2) : 0 \leq x \leq 1, \lambda_1 = 0, \lambda_2 = 0\}$$

Hence the global optimal solution for the MPCC problem is  $x = \frac{1}{2}$  with optimal value  $f(\frac{1}{2}) = 0$ .

The feasible set of the SBP problem is

$$\operatorname{argmin}\{h(x) : 0 \leq x \leq 3\} = \{x : 0 \leq x \leq 1\}$$

Therefore the global solution of the SBP problem is same as the MPCC i.e.  $x = \frac{1}{2}$ .

The SBP and MPCC problems are different in general if the Slater's condition is not satisfied. Next we present some examples to show how they are different.

### Example

SBP has unique solution but corresponding MPCC is not feasible (Slater's condition is not satisfied).

Let

$$f(x_1, x_2) = x_1 + x_2$$

$$h(x_1, x_2) = x_1$$

$$g_1(x_1, x_2) = x_1^2 - x_2$$

$$g_2(x_1, x_2) = x_2$$

Clearly,  $g_1(x_1) \leq 0$  and  $g_2(x_1, x_2) \leq 0$  together imply that  $x_1 = 0 = x_2$ . Which implies that Slater's condition fails for the lower level problem of the SBP.

Now, the feasible set for the SBP problem is

$$\operatorname{argmin}\{h(x_1, x_2) : x_1 = 0 = x_2\} = \{(0, 0)\}$$

Therefore,  $(0, 0)$  is the solution of the SBP problem.

But for  $x_1 = 0 = x_2$ , there does not exist  $\lambda_1 \geq 0$  and  $\lambda_2 \geq 0$  such that

$$\nabla h(x_1, x_2) + \lambda_1 \nabla g_1(x_1, x_2) + \lambda_2 \nabla g_2(x_1, x_2) = 0$$

Therefore the MPCC problem is not feasible even when the SBP has unique solution in case of the failure of Slater's condition.

### Example

SBP and the corresponding MPCC both are feasible but have different solution sets (Slater's condition is not satisfied).

Let

$$f(x, y) = (x - 1)^2 + y^2$$

$$h(x, y) = x^2 y$$

$$g_1(x, y) = y^2$$

$$g_2(x, y) = -x$$

Now,  $g_1(x, y) \leq 0$  and  $g_2(x, y) \leq 0$  together implies that

$$x \geq 0 \text{ and } y = 0.$$

Therefore,

$$\operatorname{argmin}\{h(x, y) : x \geq 0, y = 0\} = \{(x, 0) : x \geq 0\}$$

Hence,  $(1, 0)$  is the solution of the SBP problem with optimal value  $f(1, 0) = 0$ .

Now for the MPCC problem and  $x \geq 0, y = 0$

$$\nabla h(x, y) + \lambda_1 \nabla g_1(x, y) + \lambda_2 \nabla g_2(x, y) = 0$$

holds true if  $x = 0, y = 0, \lambda_1 \geq 0, \lambda_2 = 0$ .

Therefore,  $x = 0, y = 0$  is the optimal solution for the MPCC problem with optimal value  $f(0, 0) = 1$  which is different from the SBP problem.

# References



Dempe, S.; Dutta, J. Is bilevel programming a special case of a mathematical program with complementarity constraints? *Math. Program.* 131 (2012), no. 1-2, Ser. A, 3748.

THANK YOU