

Lecture 19: October 14

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This lecture's notes illustrate some uses of various $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X}$ macros. Take a look at this and imitate.

19.1 Multistage Games with Observed Actions (Feedback Games)

We will consider a particular class of Multi-act games, called *feedback games*.

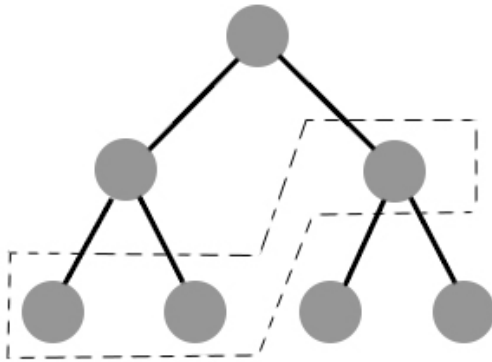
Definition 19.1.1 *A feedback game in extensive form is :*

- *A tree divided into 'stages'*
- *First acting player at each stage has singleton information set*
- *Information sets of other players are such that none of them includes nodes corresponding to branches emanating from 2 or more information sets of the first acting player*
- *No information set contains nodes from more than 1 stage*
- *Payoffs are stored at leaf nodes*

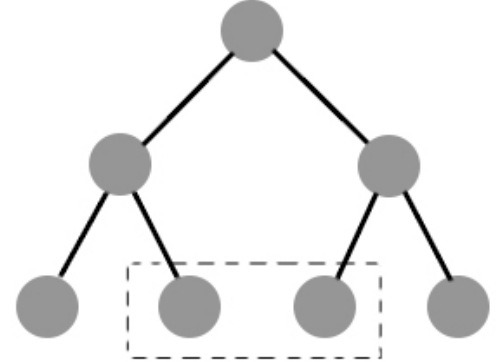
Further, we can assume that :

- *Each stage is a single-act game*
- *If single-act games at each stage are ladder nested(or nested) then it is a ladder nested(or nested) game*
- *Every player plays once in each stage*
- *Each player knows the state of the game at every level of play*

Therefore such trees are not possible :



Set contains nodes of the tree belonging to different levels of play



Set contains nodes corresponding to branches emanating from 2 different information sets of the other player

19.2 Nash Equilibrium in feedback games

Let γ^i be the strategy for player P_i , such that

$$\gamma^i = \gamma_1^i, \gamma_2^i, \dots, \gamma_k^i \tag{19.1}$$

where k is total no of stages and γ_j^i is the strategy of P_i at stage j .

Then Nash Equilibrium for the game will be,

$$J^i(\gamma^{i*}, \gamma^{-i*}) \leq J^i(\gamma^i, \gamma^{-i*}) \quad \forall \gamma^i \tag{19.2}$$

19.2.1 Feedback Nash Equilibrium

Let $\gamma^{1*}, \gamma^{2*}, \dots, \gamma^{N*}$ be such that they satisfy Nash Equilibrium condition at every stage, then such a Nash Equilibrium is called **Feedback Nash Equilibrium**.

A feedback N.E can be found using backward induction.

N.E condition at stage k ,

$$J^i(\gamma_1, \dots, \gamma_{k-1}, \gamma_k^{i*}, \gamma_k^{-i*}) \leq J^i(\gamma_1, \dots, \gamma_{k-1}, \gamma^i, \gamma_k^{-i*}) \quad \forall \gamma_k^i, \forall i, \forall \gamma_1, \dots, \gamma_{k-1} \tag{19.3}$$

where $\gamma_t = \gamma_t^1, \dots, \gamma_t^N$, i.e. strategy of all N players at stage t .

Therefore, using backward induction, N.E condition at stage k-1 would be,

$$J^i(\gamma_1, \dots, \gamma_{k-2}, \gamma_{k-1}^{i*}, \gamma_{k-1}^{-i*}, \gamma_k^*) \leq J^i(\gamma_1, \dots, \gamma_{k-2}, \gamma_{k-1}^i, \gamma_{k-1}^{-i*}, \gamma_k^*) \quad \forall \gamma_{k-1}^i, \forall i, \forall \gamma_1, \dots, \gamma_{k-2} \quad (19.4)$$

Proceeding in this manner, N.E condition at stage 1 would be,

$$J^i(\gamma_1^{i*}, \gamma_1^{-i*}, \gamma_2^*, \dots, \gamma_k^*) \leq J^i(\gamma_1^i, \gamma_1^{-i*}, \gamma_2^*, \dots, \gamma_k^*) \quad \forall \gamma_1^i, \forall i, \quad (19.5)$$

Theorem 19.2.1 *Every Feedback NE is a NE*

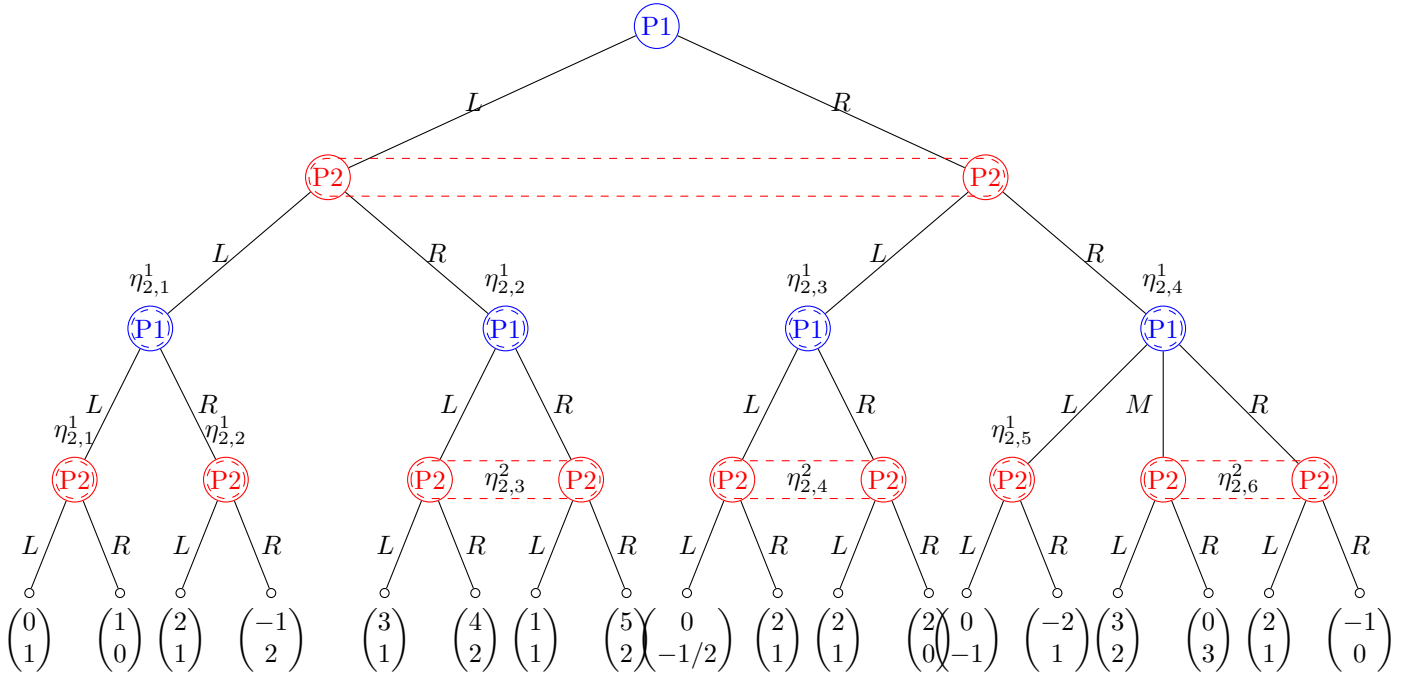
Hint: Proof by Contradiction

19.3 Algorithm for Feedback Nash Equilibrium

- Starting from stage k solve each single-act game for information set of first acting player.
- Delete subtree and replace with equilibrium payoff Work this way to the root of tree.
- If single-act game admit multiple equilibrium, repeat for each.

19.3.1 Example

Consider the following game tree in extensive form.



Lemma 19.1 *A Feedback Nash Equilibrium is of type Delayed Commitment Type if it induces a delayed commitment equilibrium in each stage.*

It can be shown that the game has 4 single-act games each starting from P1 node at level 2 in the tree. Equilibrium strategies of these 4 single-act games are as follows:

First single-act game:

$$\gamma_2^{1*}(\eta_{2,1}^1) = L \tag{19.6}$$

$$\gamma_2^{2*}(\eta_{2,1}^2) = \begin{cases} R & \text{if } U_2^1 = L \\ L & \text{Otherwise} \end{cases} \tag{19.7}$$

Second single-act game:

$$\gamma_2^{1*}(\eta_{2,2}^1) = R \tag{19.8}$$

$$\gamma_2^{2*}(\eta_{2,3}^2) = L \tag{19.9}$$

Third single-act game:

$$\gamma_2^{1*}(\eta_{2,3}^1) = L \tag{19.10}$$

$$\gamma_2^{2*}(\eta_{2,4}^2) = L \tag{19.11}$$

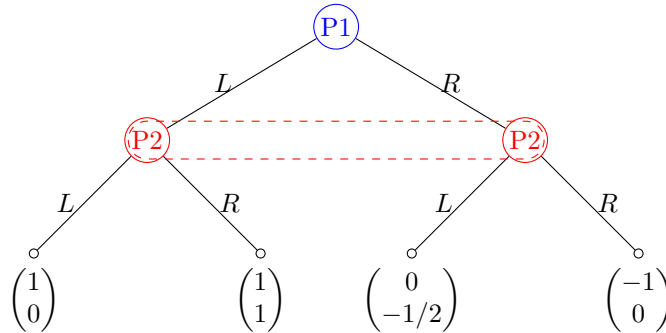
Fourth single-act game: This has two equilibria. One from the left subtree and second from the right subtree. From the right sub-tree we get static equilibrium where as from left sub-tree we get equilibrium of delayed commitment type. Let us analyze them one by one.

Case1: Static equilibrium

$$\gamma_2^{1*}(\eta_{2,4}^1) = R \tag{19.12}$$

$$\gamma_2^{2*}(\eta_{2,4}^2) = L \tag{19.13}$$

Now following step 2 of the algorithm mentioned above we can reduce the tree to the following subtree:



The equilibrium of the game is:

$$\gamma_1^{1*} = R \tag{19.14}$$

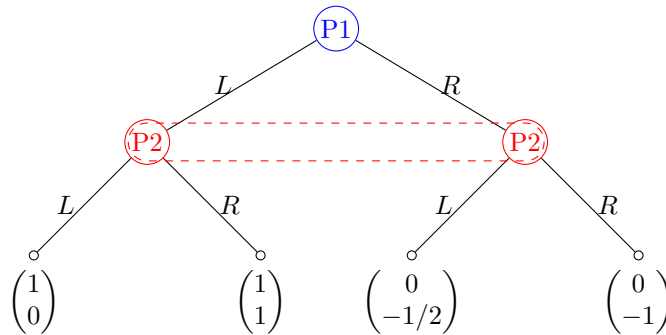
$$\gamma_1^{2*} = L \tag{19.15}$$

Case2: Delayed Commitment Type Equilibrium

$$\gamma_2^{1*}(\eta_{2,4}^1) = L \tag{19.16}$$

$$\gamma_2^{2*}(\eta_{2,5}^2) = \begin{cases} L & \text{if } U_2^1 = L \\ R & \text{Otherwise} \end{cases} \quad (19.17)$$

Now following step 2 of the algorithm mentioned above we can reduce the tree to the following subtree:



The equilibrium of the game is:

$$\gamma_1^{1*} = R \quad (19.18)$$

$$\gamma_1^{2*} = R \quad (19.19)$$

Note: This is a Feedback Equilibrium of Delayed Commitment Type.