

# Indian Institute of Technology Bombay Department of Electrical Engineering

EE-309: Microprocessors

## **Assignment 2**

Design a 6 stage pipelined processor, IITB-RISC14, whose instruction set architecture is provided. *IITB-RISC14* is a 16-bit very simple computer developed for the teaching that is based on the Little Computer Architecture. The *IITB-RISC14* is an 8-register, 16-bit computer system. It should follow the standard 6 stage pipelines (Instruction fetch, instruction decode, register read, execute, memory access, and write back). The architecture should be optimized for performance, i.e., should include hazard mitigation techniques. Hence, it should have forwarding and branch prediction technique.

Group: Group of THREE

Submission deadlines:

Nov 09 (Sunday):

Full design description with forwarding and branch predictor (on paper).

Nov 25 (Tuesday)

Verilog/VHDL code of the design (without branch predictor). Simulation results and testbench.

#### **IITB-RISC14 Instruction Set Architecture**

*IITB-RISC14* is a 16-bit very simple computer developed for the teaching that is based on the Little Computer Architecture. The *IITB-RISC14* is an 8-register, 16-bit computer system. It has 8 general-purpose registers (R0 to R7). Register R7 is always stores Program Counter. All addresses are short word addresses (i.e. address 0 corresponds to the first two bytes of main memory, address 1 corresponds to the second two bytes of main memory, etc.). This architecture uses condition code register which has two flags Carry flag © and Zero flag (Z). The *IITB-RISC14* is very simple, but it is general enough to solve complex problems. The architecture allows predicated instruction execution and multiple load and store execution. There are three machine-code instruction formats (R, I, and J type) and a total of 14 instructions. They are illustrated in the figure below.

#### **R** Type Instruction format

Opcode	Register A (RA)	Register B (RB)	Register B (RB)	Unused	Condition (CZ)
(4 bit)	(3 bit)	(3-bit)	(3-bit)	(1 bit)	(2 bit)

#### I Type Instruction format

Opcode	Register A (RA)	Register C (RC)	Immediate
(4 bit)	(3 bit)	(3-bit)	(6 bits signed)

#### J Type Instruction format

Opcode	Register A (RA)	Immediate
(4 bit)	(3 bit)	(9 bits signed)

### **Instructions Encoding:**

instructions Encoding:						
ADD:	00_00	RA	RB	RC	0	00
ADe:	00_00	RA	RB	RC	0	10
ADZ:	00_00	RA	RB	RC	0	01
ADI:	00_01	RA	RB		6 bit Immediat	e
NDI.	00_10	RA	RB	RC	0	00
NDC:	00_10	RA	RB	RC	0	10
NDZ:	00_10	RA	RB	RC	0	01
LHI	00_11	RA		9 bit Immediate		
LW:	01_00	RA	RB	(	6 bit Immediat	e
SW.	01_01	RA	RB	6 bit Immediate		
LM:	01_10	RA	0 + 8	- 8 bits corresponding to Reg R7 to R0		
SM:	01_11	RA	0 + 8 bits corresponding to Reg R7 to R0			to R0
BEQ.	11_00	RA	RB		6 bit Immediat	е
JAL:	10_00	RA	9 bit Immediate offset			
JLR:	10_01	RA	RB		000_000	
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RA: Register A

RB: Register B

RC: Register C

## **Instruction Description**

	Mnemonic	Name & Format	Assembly	Action
	ADD	ADD (R)	Add rc, ra, rb	Add content of regB to regA and store result in regC.
	ADC	Add if carry set (R)	Adc rc, ra, rb	Add content of regB to regA and store result in regC, if carry flaf is set.  It sets C & Z flags
_	ADZ	Add if zero set (R)	Adz rc, ra, rb	Add content of regB to regA and store result in regC, if zero flag is set.  It sets C & Z flags
	ADL	Add immediate (I)	Adi rb, ra, imm6	Add content of regA with Imm (sign extended) and store result in regB.  It set C and Z flags
_	NDU	Nand (R)	Ndu rc, ra, rb	NAND the content of regB to regA and store result in regC.
	NDC	Nand if carry set (R)	Ndc rc, ra, rb	NAND the content of regB to regA and store result in regC if carry flag is set.
	NDZ	Nand if zero set (R)	Ndc rc, ra, rb	NAND the content of regB to regA and store result in regC if zero flag is set.
	LHI	Load higher immediate (J)	Lhi ra, Imm	Place 9 bits immediate into most significant 9 bits of register A (RA) and lower 7 bits are assigned to zero.
	LW	Load (I)	Lw ra, rb, Imm	Load value from memory into reg A.  Memory address is formed by adding immediate 6 bits with content of red B.

				It sets zero flag.
_	SW	Store (I)	Sw ra, rb, Imm	Store value from reg A into memory.  Memory address is formed by adding immediate 6 bits with content of red B.
	LM	Load multiple (J)	Lw ra, Imm	Load multiple registers whose address is given in the immediate field (one bit per register, R7 to R0) in order from right to left, i.e, registers from R0 to R7 if corresponding bit is set. Memory address is given in reg A. Registers are loaded from consecutive addresses.
	SM	Store multiple (J)	Sm, ra, Imm	Store multiple registers whose address is given in the immediate field (one bit per register, R7 to R0) in order from right to left, i.e, registers from R0 to R7 if corresponding bit is set. Memory address is given in reg A. Registers are stored to consecutive addresses.
_	BEQ	Branch on Equality (I)	Beq ra, rb, Imm	If content of reg A and regB are the same, branch to PC+Imm, where PC is the address of beq instruction
_	JAL	Jump and Link (I)	Jalr ra, Imm	Branch to the address PC+ Imm.  Store PC+1 into regA, where PC is the address of the jalr instruction
-	JLR	Jump and Link to Register (I)	Jalr ra, rb	Branch to the address in regB.  Store PC+1 into regA, where PC is the address of the jalr instruction